

# Solution Extraction from Long-distance Resolution Proofs

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**International Workshop on Quantified Boolean Formulas**



FAKULTÄT  
FÜR INFORMATIK

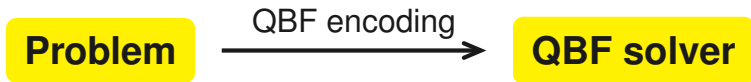
Faculty of Informatics



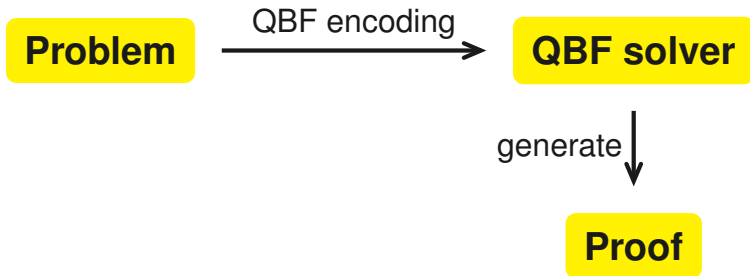
# Overview

**Problem**

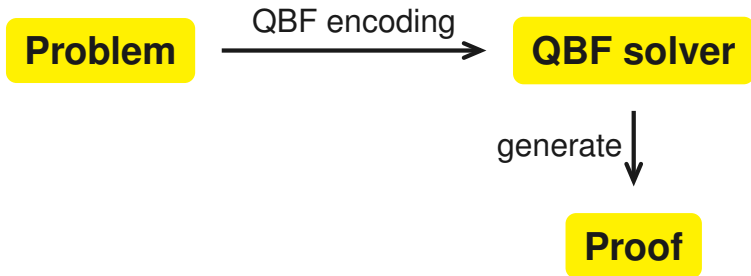
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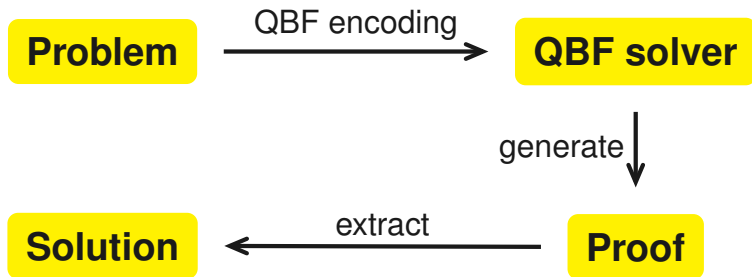


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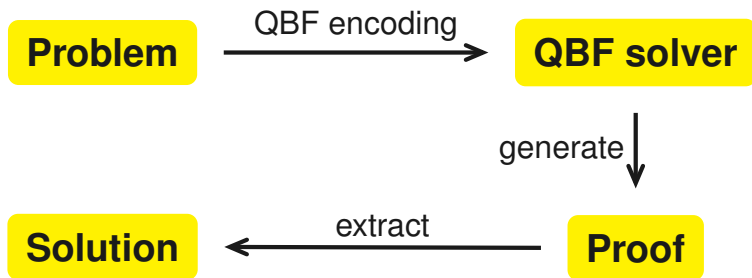
E.g. Q-Resolution by  
DepQBF  
[Lonsing and Biere, 2010]

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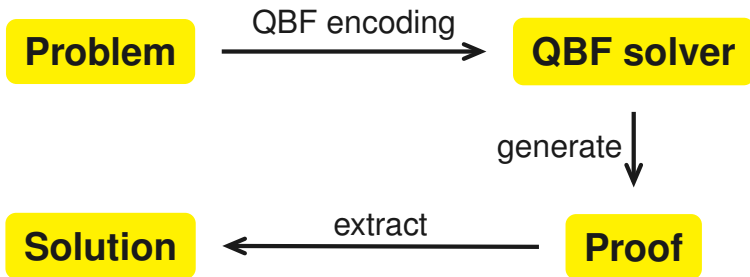
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Winning strategy for  $\forall$  player  
[Goultiaeva et al., 2011]

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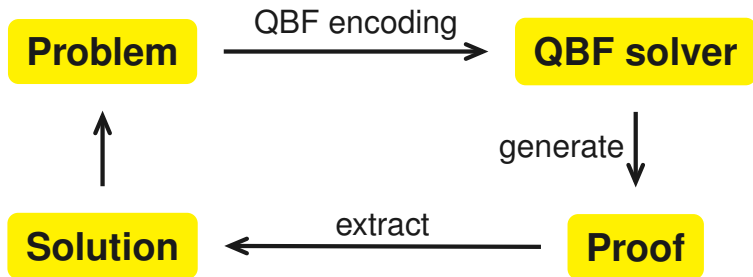
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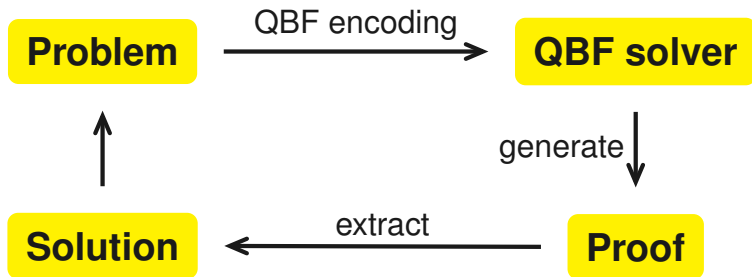


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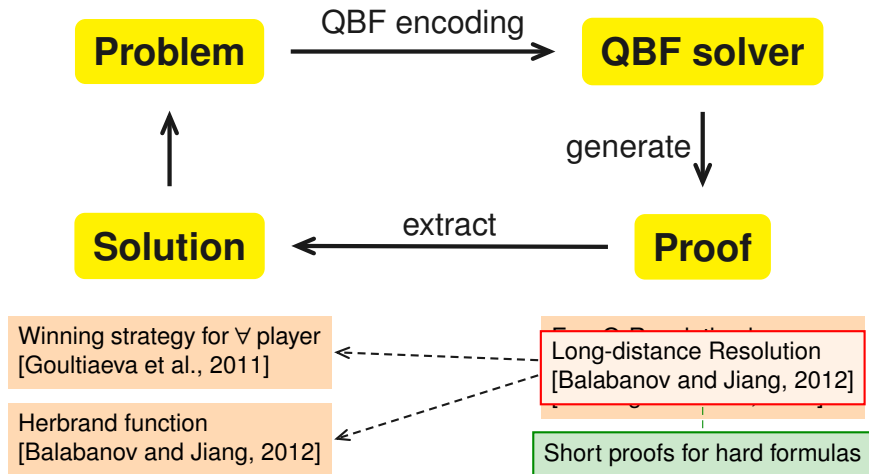


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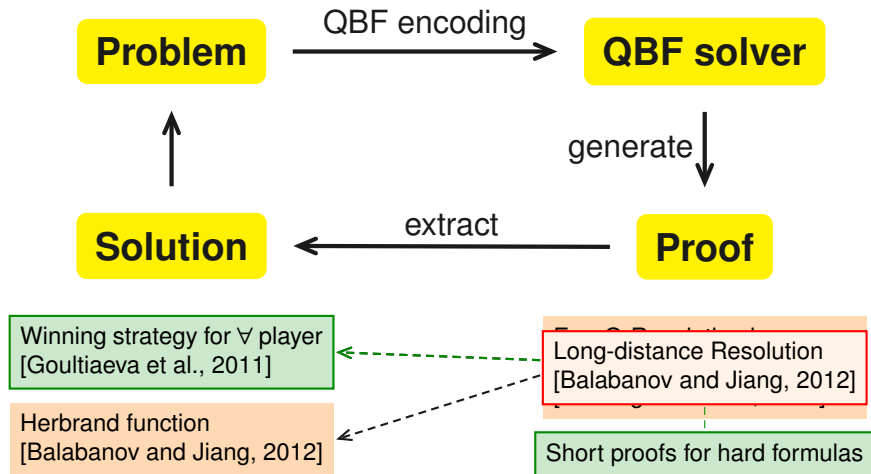
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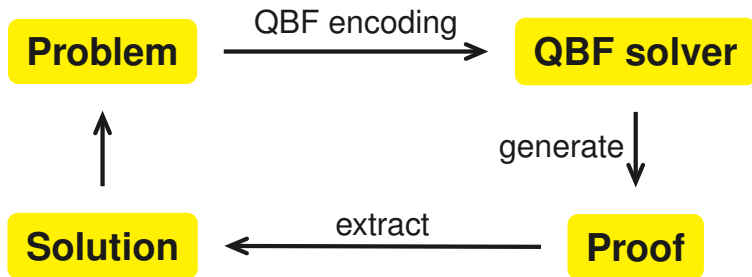
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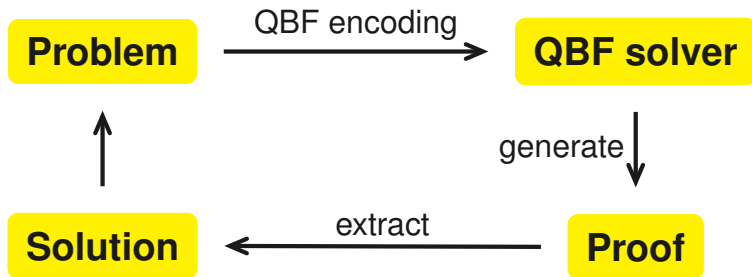


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# Solution Representations

## Winning strategy [Goultiaeva et al., 2011]

- ▶ Game-theoretic view
- ▶ Controller computes move for the  $\forall$  player according to previous moves of the  $\exists$  player
- ▶ Runtime polynomial *in the size of the refutation*

# Solution Representations

## Winning strategy [Goultiaeva et al., 2011]

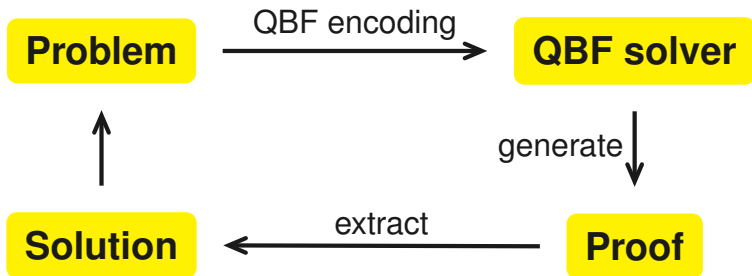
- ▶ Game-theoretic view
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## Herbrand functions [Balabanov and Jiang, 2012]

- ▶ Function for each  $\forall$  variable depending on  $\exists$  variables to its left
- ▶ Replacing each  $\forall$  variable by its function renders the matrix unsatisfiable
- ▶ Runtime polynomial *in the size of the refutation*



# Overview

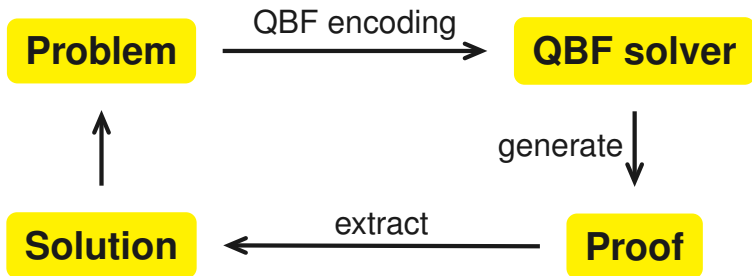


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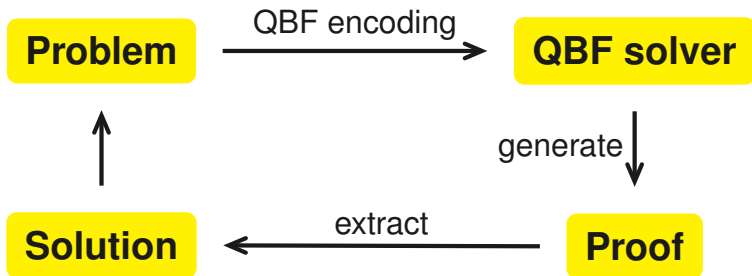
Herbrand function  
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*polynomial*

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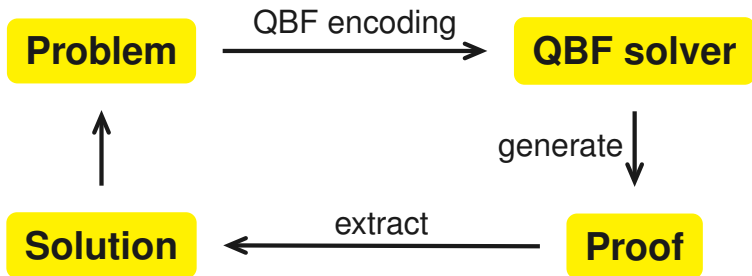
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**exponential**

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Long-distance Resolution  
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# Long-distance Resolution (LDQ-Resolution)

Q-resolution with the following additional derivation rules

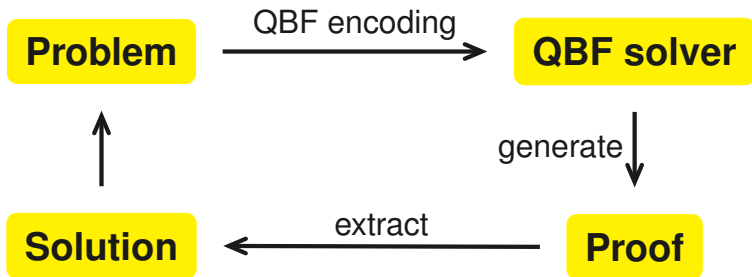
$$\frac{K^l \vee p \vee x \quad K^r \vee \bar{p} \vee \bar{x}}{K^l \vee K^r \vee x^*} \quad \text{lev}(p) < \text{lev}(x)$$

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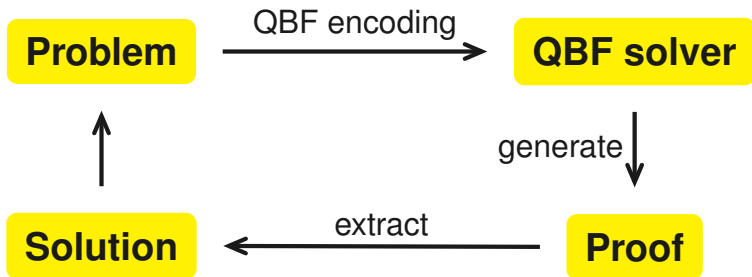


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Short proofs for hard formulas

# Short Proofs for Hard Formulas

Family  $(\varphi_t)_{t \geq 1}$  of QBFs in PCNF with prefix

$$\exists d_0 d_1 e_1 \forall x_1 \exists d_2 e_2 \forall x_2 \exists d_3 e_3 \cdots \forall x_{t-1} \exists d_t e_t \forall x_t \exists f_1 \cdots f_t$$

and matrix:

$K_0$	$\overline{d_0}$	$K_1$	$d_0 \vee \overline{d_1} \vee \overline{e_1}$	
$K_{2j}$	$d_j \vee \overline{x_j} \vee \overline{d_{j+1}} \vee \overline{e_{j+1}}$	$K_{2j+1}$	$e_j \vee x_j \vee \overline{d_{j+1}} \vee \overline{e_{j+1}}$	$j = 1, \dots, t-1$
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- ▶ Any Q-refutation for  $\varphi_t$  is exponential. [Kleine Büning et al., 1995]
- ▶ Q-resolution plus resolution over  $\forall$  variables yields polynomial refutations. [Van Gelder, 2012]

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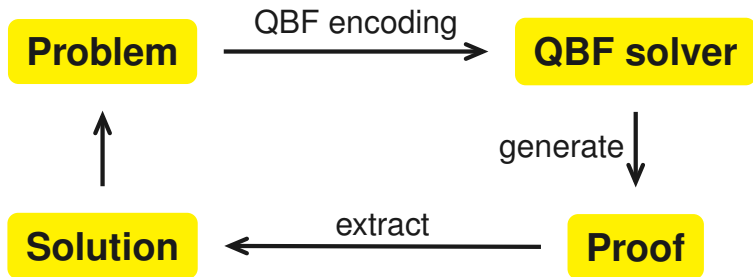
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- ▶ Q-resolution plus resolution over  $\forall$  variables yields polynomial refutations. [Van Gelder, 2012]
- ▶ Polynomial LDQ-refutation is possible. [This work]

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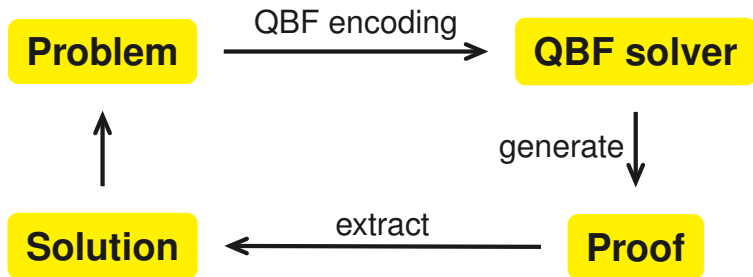
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Short proofs for hard formulas

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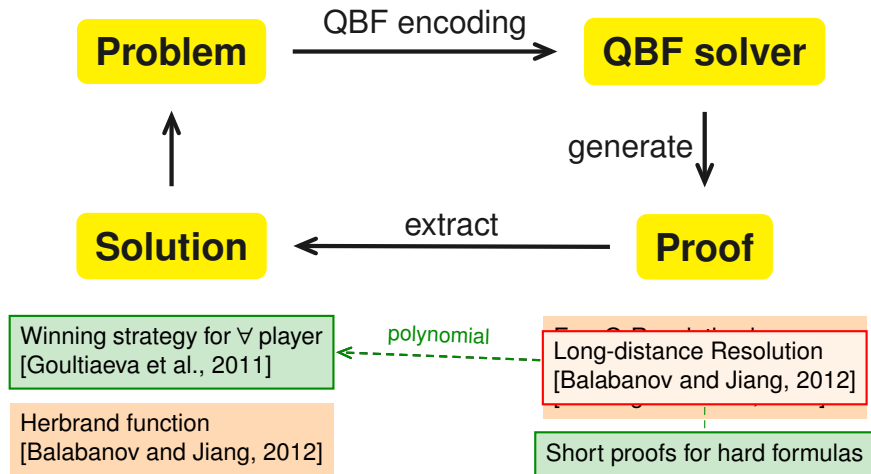
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Short proofs for hard formulas

← ?

# Overview



# Strategy Extraction from LDQ-Refutations

$\exists a \forall x \exists b \forall y \exists c$

$(a, x, b, y, c)$

$(\bar{a}, \bar{y})$

$(\bar{c})$

$(\bar{b})$

# Strategy Extraction from LDQ-Refutations

LDQ-Resolution

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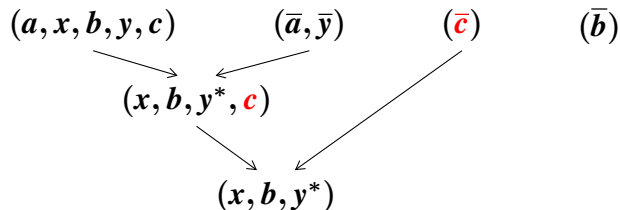
$(x, b, y^*, c)$



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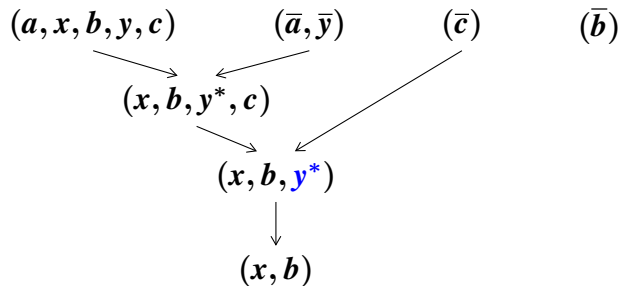
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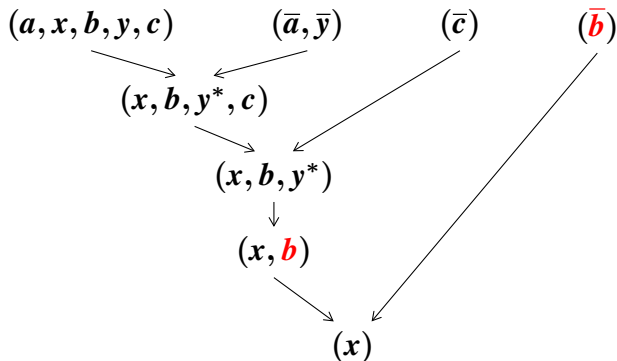
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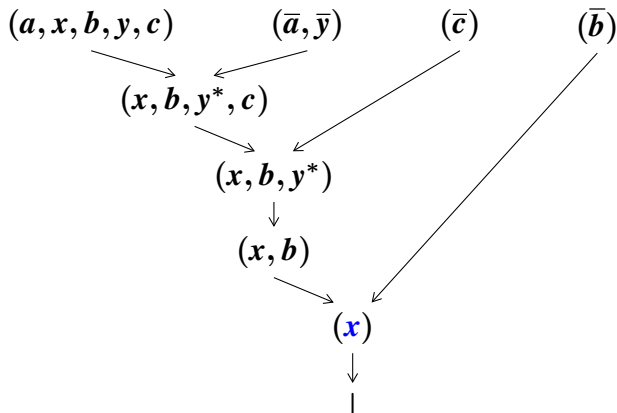
LDQ-Resolution



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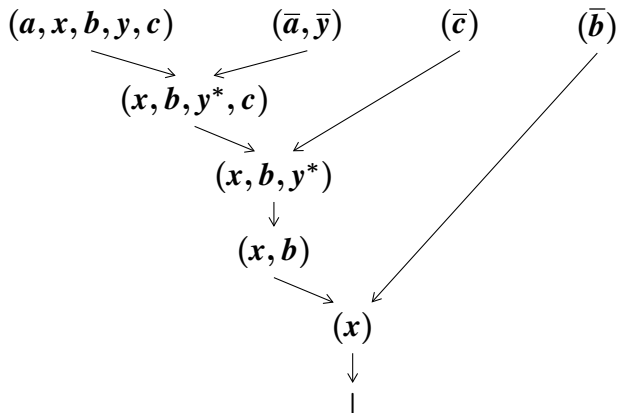
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LDQ-Resolution



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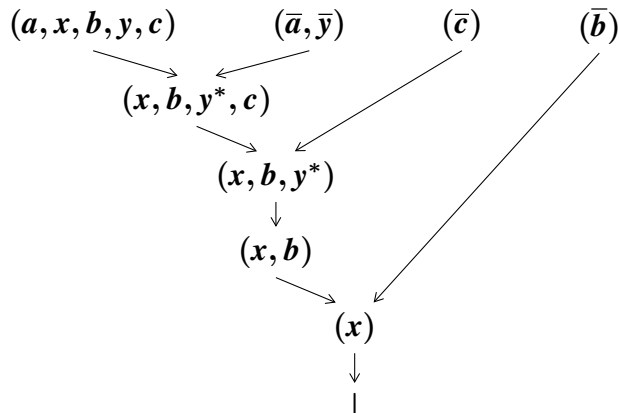
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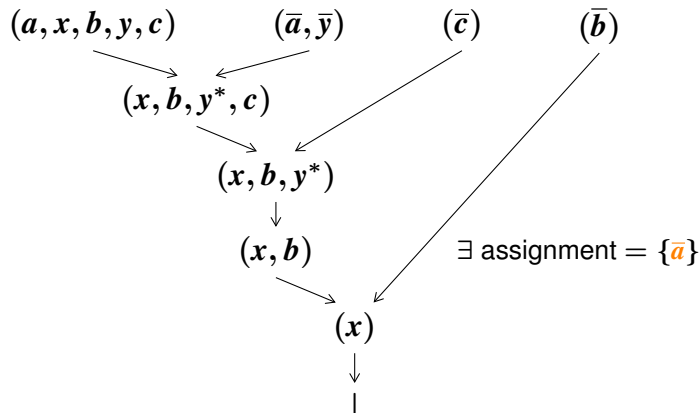
Strategy Extraction



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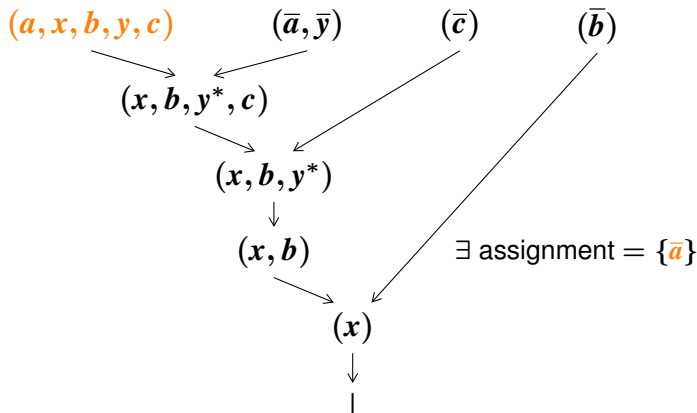
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Strategy Extraction



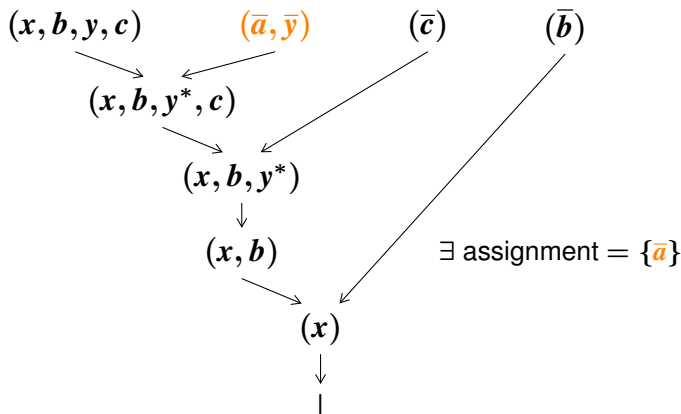




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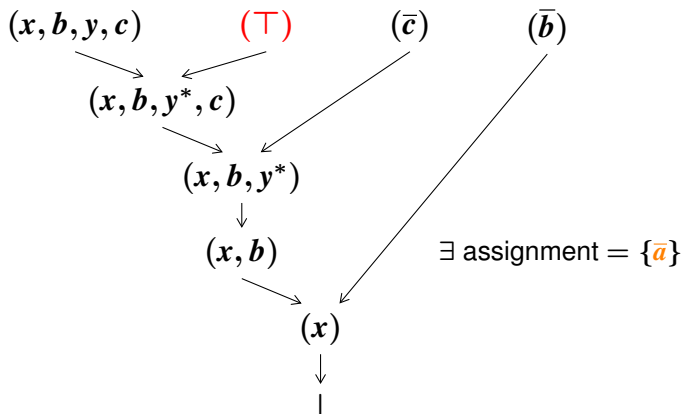
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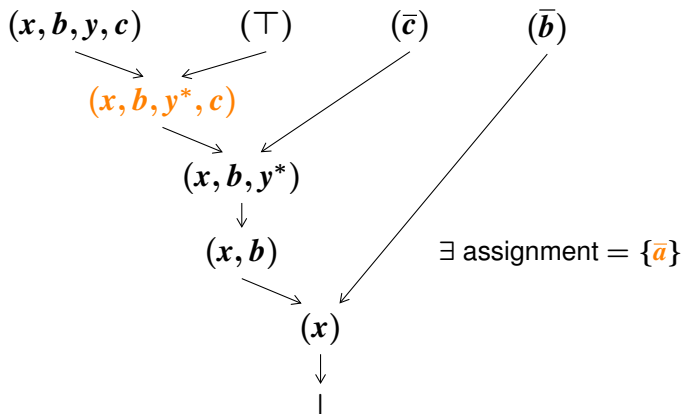
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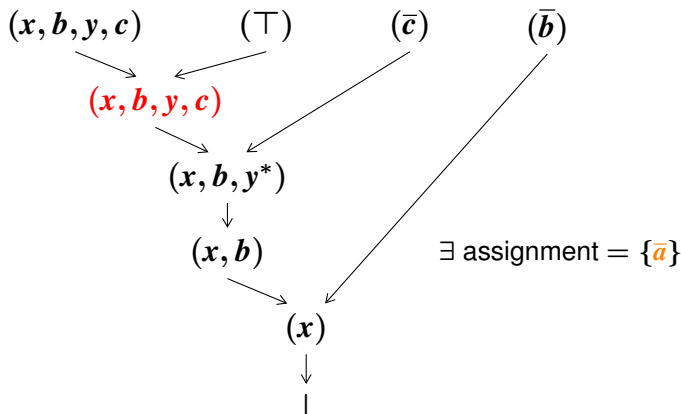
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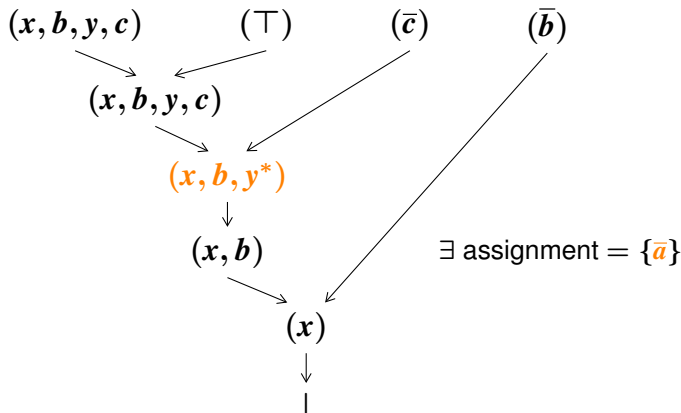
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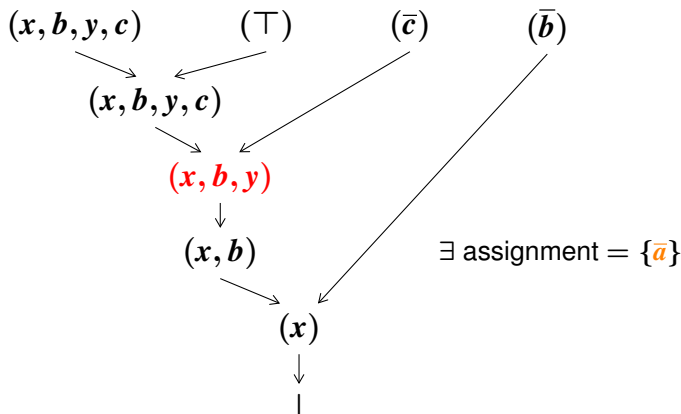
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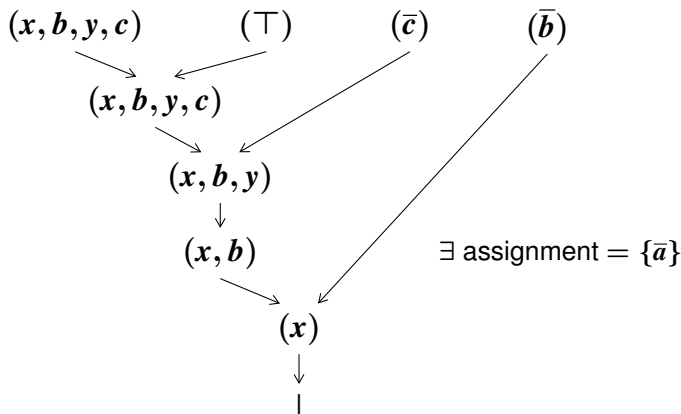
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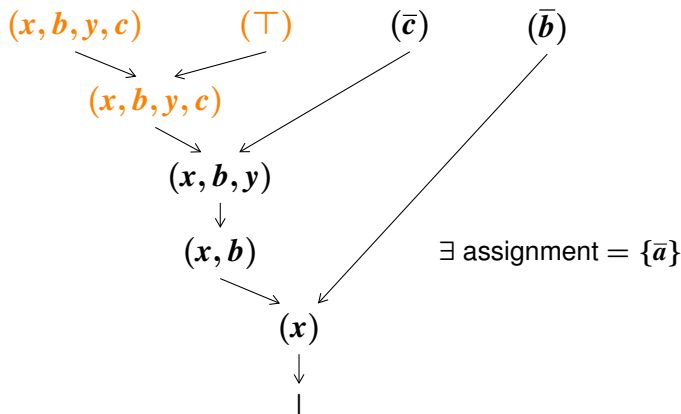




# Strategy Extraction from LDQ-Refutations

$\exists a \forall x \exists b \forall y \exists c$

Strategy Extraction



# Strategy Extraction from LDQ-Refutations

$\forall x \exists b \forall y \exists c$

Strategy Extraction

$(x, b, y, c)$

$(\bar{c})$

$(\bar{b})$

$(x, b, y)$

$(x, b)$

$\exists$  assignment =  $\{\bar{a}\}$

$(x)$

$\perp$

# Strategy Extraction from LDQ-Refutations

Strategy Extraction

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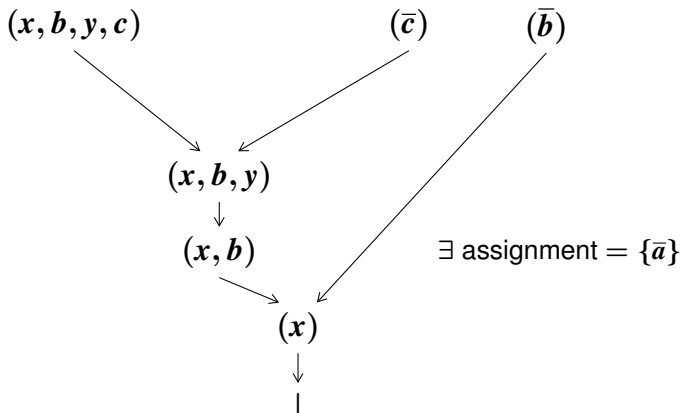
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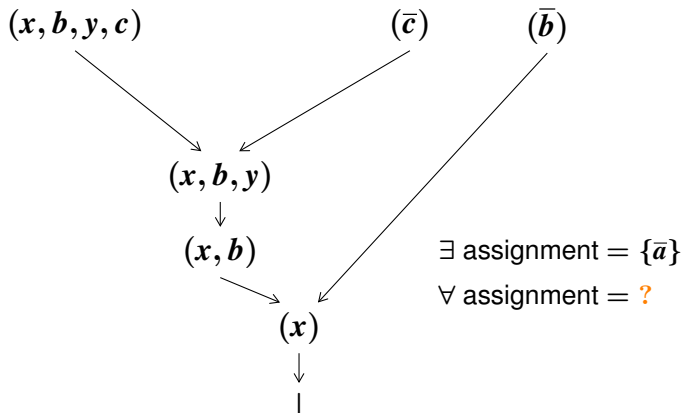
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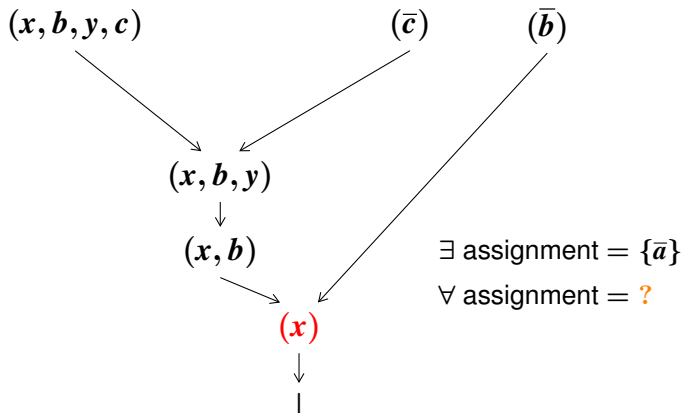
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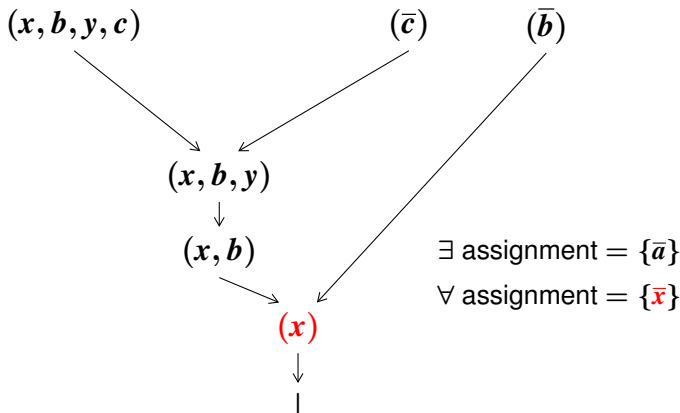
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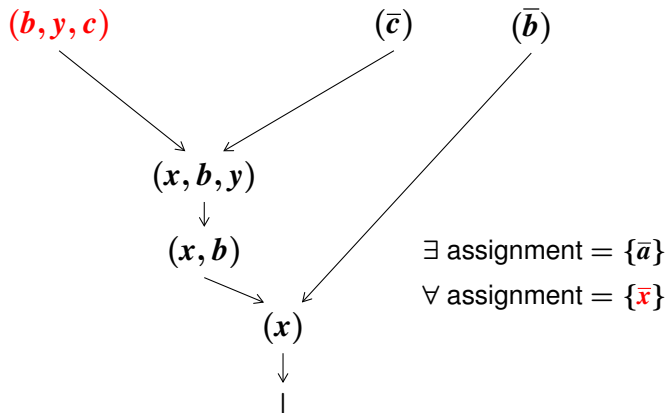
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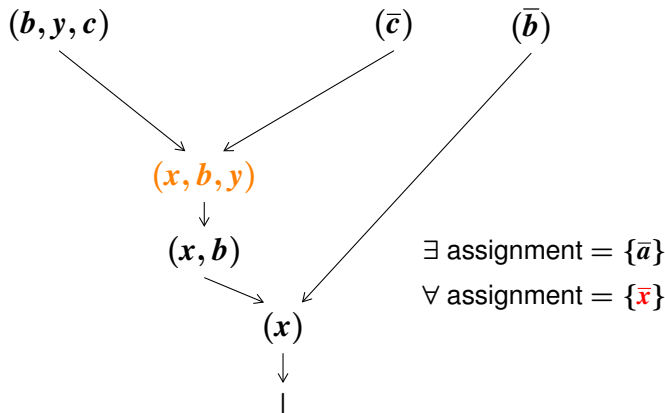
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# Strategy Extraction from LDQ-Refutations

$\forall x \exists b \forall y \exists c$

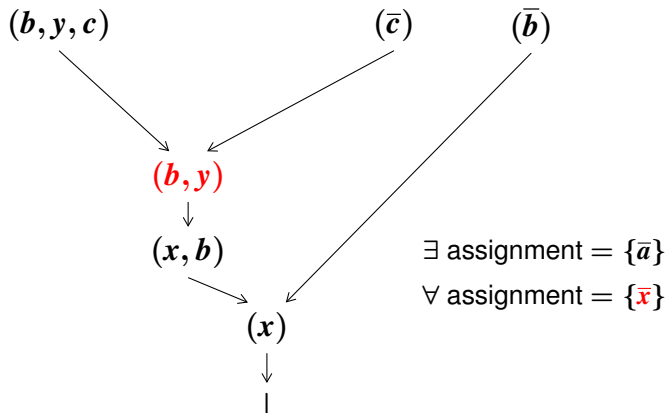
Strategy Extraction



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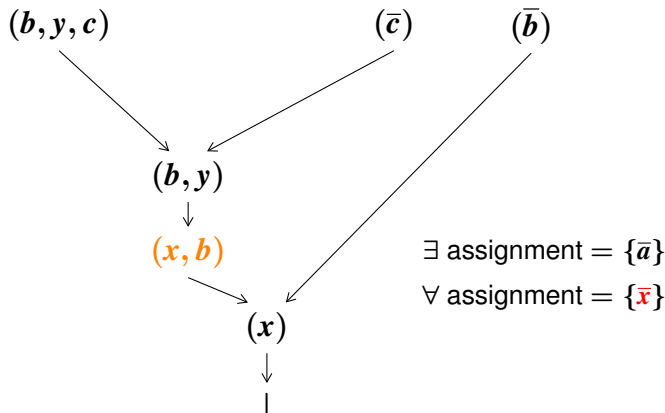
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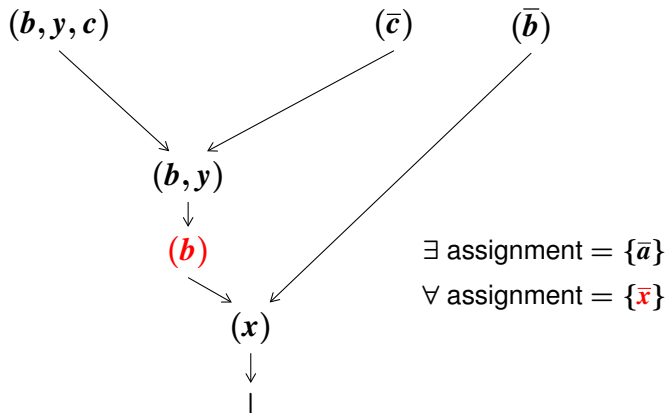
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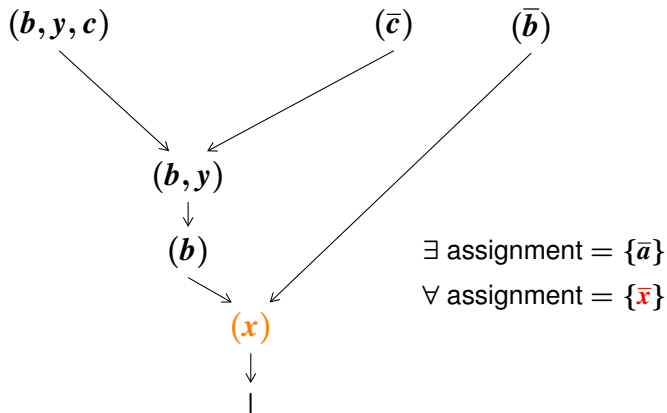
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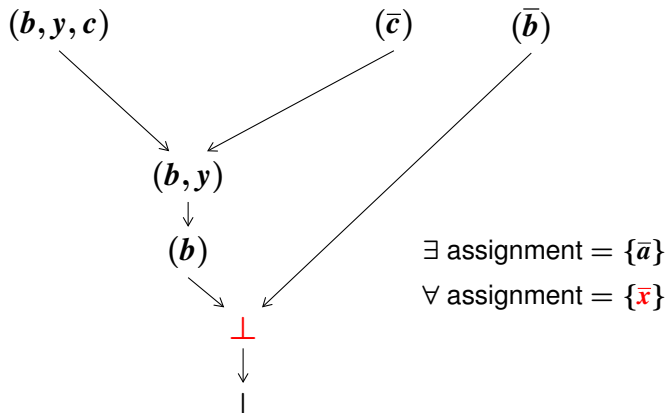
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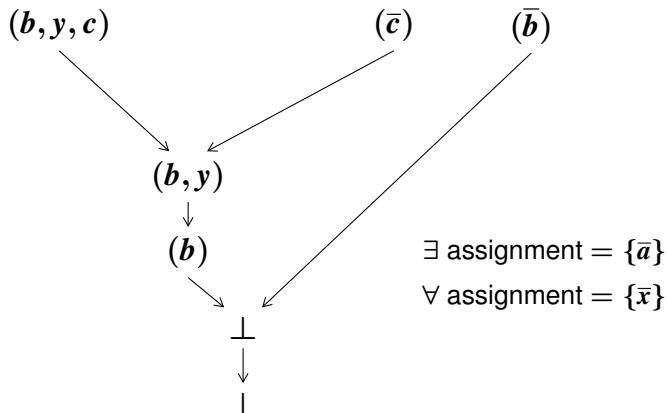
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Strategy Extraction

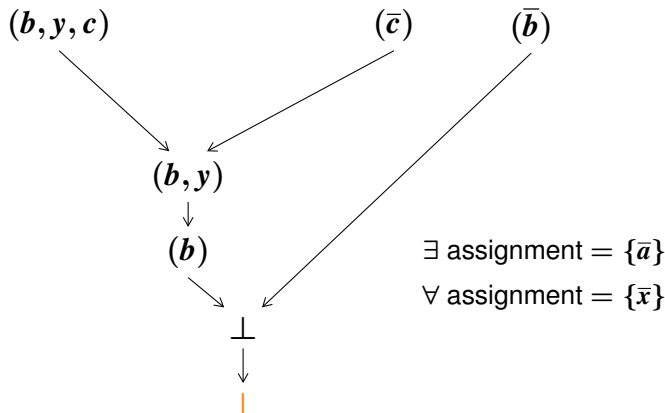




# Strategy Extraction from LDQ-Refutations

$\forall x \exists b \forall y \exists c$

Strategy Extraction



# Strategy Extraction from LDQ-Refutations

Strategy Extraction

$\exists b \forall y \exists c$

$(b, y, c)$

$(\bar{c})$

$(\bar{b})$

$(b, y)$

$(b)$

$\perp$

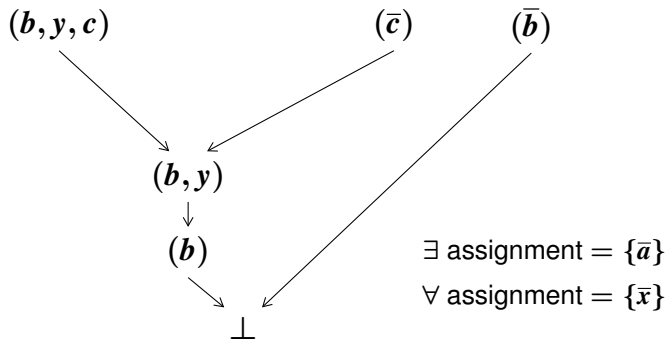
$\exists$  assignment =  $\{\bar{a}\}$

$\forall$  assignment =  $\{\bar{x}\}$

# Strategy Extraction from LDQ-Refutations

$\exists b \forall y \exists c$

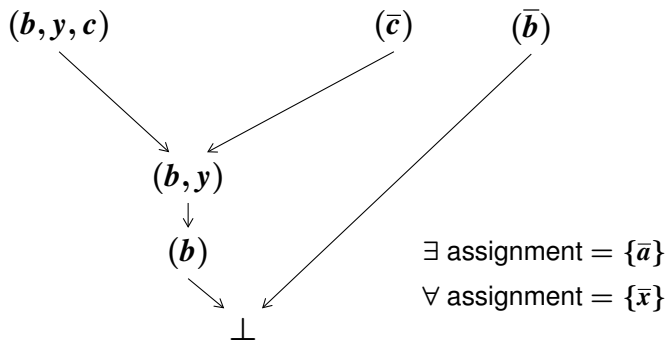
Strategy Extraction



# Strategy Extraction from LDQ-Refutations

$\exists b \forall y \exists c$

Strategy Extraction



... and so on

# Conclusions and Future Work

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- ▶ LDQ-resolution support in QBF solvers

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